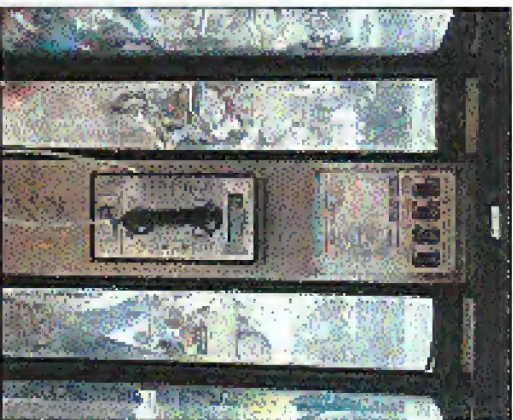


BACK COVER Foreign Phones



Cambodia. A card phone in a busy street in Banteay Meanchey.

Photo by Eric Tucker



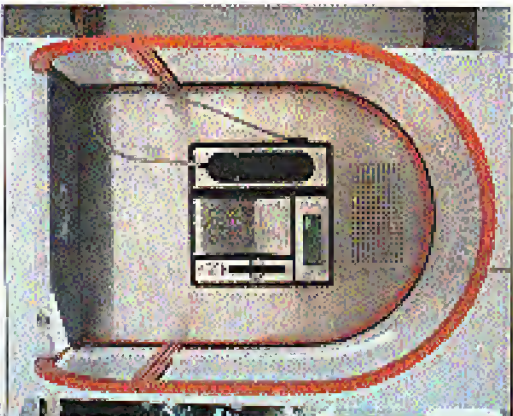
Cambodia. Another card phone from the capital city of Phnom Penh. It's rumored that there are no coin phones at all in the country.

Photo by Eric Tucker



Greece. Found in the small Greek village of Milos.

Photo by John Klausmann



Greece. From the village of Messinias. We hope that isn't a spoiler behind the grill above the phone since it looks like it would easily disable anyone standing it see.

Photo by John Klausmann

Look on the other side of this page for even more photos!

26000

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"We all have to fight against the hacker community." - Judy Elder of Microsoft

Canada, as quoted by the CBC, July 31, 2001



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Shout Outs: the people who came together to make HAL 2001 work, those who continue to help us all get through a period of unimaginable darkness in NYC

RIP WTC

Dedicated to the memory of Wau Holland (11/20/1951-07/29/2001) and the thousands lost in New York on September 11

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Consequences

It takes an event of great magnitude to really put things into perspective, to make us realize how insignificant our daily concerns can be. At the same time, such an occurrence can trigger a chain of events that wind up magnifying these concerns.

It's hard to imagine anyone who hasn't felt the horrible weight of September 11. Here, before our eyes, was all the confirmation we needed to see how uncalibrated the human race could be and how vulnerable we as individuals and a society really are to those who value neither.

We feel the outrage along with everyone else. Anyone responsible for such heinous acts, whether directly or by helping to organize them, deserves no mercy from any court in the world.

Rage, however, often makes us lose sight of some of the important things that we're supposed to be defending in the first place. And we have to be extremely careful not to add additional loss of freedom to the loss of life that is the legacy of terrorism.

What perhaps is most disturbing is the speed with which things began to change after the attacks. It was as if members of Congress and lawmakers were poised to spring into action the moment public opinion began to turn and before our nation's sense had a chance of regaining its dominance. Within hours of the horrific events, new restrictions on everything from employment to community along with broad new powers in law enforcement were being proposed—all with little or no public support from the terrified public.

We find it absolutely unconscionable that anyone would use such a tragedy to further their own agenda—whether it be by selling a product or pushing a wish list of legislation. We've witnessed a good amount of both recently and it's all pretty repugnant. Almost every new law that's been proposed is something we've already seen in the past—and rejected. And there is very little contained within them that would have been helpful in preventing the terrorist attacks in the first place.

Our concerns can best be summarized up by this quote: "Maybe the Senate wants to just go

ahead and adopt new abilities to wiretap our citizens. Maybe they want to adopt new abilities to go into people's computers. Maybe that will make us feel safer. Maybe. And maybe what the terrorists have done made us a little bit less safe. Maybe they have increased Big Brother in this country. If that is what the Senate wants, we can vote for it. But do we really show respect to the American people by slipping something in together, something that nobody on the floor can explain, and say we are changing the duties of the Attorney General, the Director of the CIA, the U.S. attorneys, we are going to change your rights as Americans, your rights to privacy? We are going to do it with no hearings, no debate. We are going to do it with minutes on a page that nobody can understand."

Those remarks came from Senator Patrick Leahy of Vermont, one of the few who seem to actually comprehend the serious risks we're facing. And when a senator expresses these kinds of fears, it's a good idea to pay attention. The consequences of not thinking this through are so great that they're difficult to even grasp.

We've faced some serious threats to freedom before all of this, as anyone who reads *2600* would know. This column was originally focused primarily on the case of Dmitry Sklyarov, the Russian programmer pictured on our cover with his son. (Our cover, incidentally, was designed well before the events of September 11 so the combination of the New York City skyline and halibut legs of doom is a rather sad coincidence.) As has already been widely reported, Sklyarov was arrested after giving a lecture at the Defcon conference in Las Vegas that July. The Russian company he worked for (Eloomssoft) manufactured a program called *SoftView* (eBook Processor (A/B/E/P)) which basically allowed users of Adobe's eBook Reader to translate files to Portable Document Format (PDF). Even though the software only works on legitimately purchased eBooks, our industry brethren were consider such a translation to be a violation of the Digital Millennium Copyright Act. Sklyarov, who had planned on returning home to Russia, was imprisoned for three weeks before finally being released on a

\$50,000 cash bail. Both he and his company have been charged with violating the DMCA, an offense which could land him in jail for 25 years and bankrupt the company. He is now stuck in the United States awaiting trial.

Ever since we became the first defendants to be charged with violating the DMCA last year with the DocSS case, we know that it would only be a matter of time before the arena changed from a civil court to a criminal court. (At press time, we were still awaiting the results of our appeal.) Now we've crossed over into a very ominous sort of scenario: Someone has essentially been imprisoned for figuring out how to translate one format of code into another. An American court seeks to put a foreign company out of business for being part of this effort. And despite the fact that Adobe themselves have changed their minds about pressing charges, the United States government intends to go forward with this case and many others. Leading the charge back in July was U.S. Attorney Robert S. Mueller III of San Francisco. Today he is the head of the FBI.

Before any of the really bad stuff started to happen, we were already asking ourselves: if things could possibly get any worse. If almost seems as if there is no limit as to how bad it can get.

In a strange confluence to this theme of despair, we had the heavy and optimism of HAL 2001. For all too brief a period, we could forget the worries back home and take part in what may have been the best hacker conference so far, where people from all over the world built the equivalent of a small city in the fields of the Netherlands.

It's interesting to know that such an event is still possible and, as usual, it took the Dutch to remind us of this. It is still possible for people of all cultures to come together and share everything from ideas to technology to the physical labor needed to bring it all together. And all of this in an environment where not a single security guard was seen, where the community of several thousand people took care of themselves—where, for, if any, didn't feel inspired by what the hacker community could accomplish if only given the chance.

If anything is to get us through the dark days ahead it has to be this spirit of HAL, which is really the original spirit of hackers everywhere—enthusiasm, exploration, exchange of ideas in a free and open setting. It will be quite a challenge to keep this spirit alive when there is so

much pressure to move in the other direction. But we have to and for the same reason that we resist corruption—we cannot let that which we believe in be corrupted and subverted by those who don't understand.

And they truly don't understand. As we go to press, the Anti-Terrorism Act is getting ready to be voted on without any public input. A little noticed provision would actually categorize violations of the Computer Fraud and Abuse Act as "federal terrorism offenses." It basically means that hacking offenses of all sorts (even those committed decades ago) could result in a life sentence without any hope of release. To categorize someone who lacks a web page or responses onto a computer system in the same way as someone who blows up buildings and sabotages airplanes is so outrageous as to be extremely offensive to anyone who has been a victim of the terrorism. It's hard to believe our government could be this ignorant. What's even scarier is the possibility that they know exactly where they're going on this. But ignorant or not, they cannot be allowed to continue down this path.

In some ways we are fortunate. The increasing activism of the hacker community over the years has put us in a position where we know what to do and can do it quicker and with more people than ever before. For instance, the Free Identity movement was in full swing within days after his arrest. Demonstrations occurred in multiple cities throughout the world. And public pressure was what got Adobe to back down, even though that action had no bearing on the case. Organizations like the Electronic Frontier Foundation are more alert than ever when it comes to cases that will decide the future of technology. Again, we encourage you to donate to them (www.eff.org or 1-877-454-5101) or, San Francisco CA 94110 (USA), to visit our site or www.freeid.org for updates, and to keep your eyes open on all levels for the ongoing dangers to freedom. Otherwise, we will all pay a very heavy price.

We lost some intellectual pillars and a whole lot of innocent lives on September 11. Now the pillars of freedom and justice which remain must be saved from destruction as well.



by Avdius

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Hacking Fortres on a properly configured machine is all but impossible. Hacking Fortres on a poorly configured machine is incredibly easy. Hacking Fortres on any machine inoperatively is easy.

If you just want to break Fortres, stop reading this and do the standard "Reformat boot disk-High system files-reboot, hard-disk system files-reboot" trick. And the thinking of a good excuse when the technician or a teacher comes over and asks you what the hell you are doing. If you want to understand how this pretty cool program works, then read on.

This article refers to Fortres 1.01 version 4 for Win 9x. It's the flagship product of Fortres Guard Corp (www.fortres.com). This version also runs on NT/2000. However the rest (aside from 1) included Fortres on some Win98, so everything here refers to Win 9x unless said otherwise. Most systems you will find running Fortres will be low-end Pentiums used in Ethernet that will have Win 9x, Netware, and maybe some anti-virus stuff. The good thing is they may have a government Internet restriction through a network.

First of all, I'm going to discuss Fortres security: how it loads, how it works. Next I'll tell you how to affect Fortres so you can run your programs, but still have it productive from script kiddies who want to change the host screen. Finally, I'll mention some of the weird parts of Fortres and how they could be exploited.

Fortres 1.01 simply adds a layer to Windows that checks every action you try to do against a checklist of approved actions. That's it, very simple. There is no way to break this security layer since it is loaded. If an action wasn't allowed, you will not be able to do it. Various methods exist for getting around copying or deleting files, to running certain programs, altering hosts, and more. There are two ways you can hack Fortres; you can prevent the security layer from loading (nearly impossible without drawing attention to yourself), or get into the

privilege setup program and alter the settings. Since the core of Fortres is so simple, Fortres usually consists of subprograms to prevent people from skipping this security layer from being loaded. Fortres also uses lots and takes files to hide what files truly do what. In fact, even in the Fortres 1.01 help file they lie to people who have legally purchased the software.

To prevent the loading process, Fortres modifies MSDOS.SYS, AUTOEXEC.BAT, and CONFIG.SYS. It makes backups of the old files, renaming them with the .DWF extension. MSDOS.SYS is appended with the following: Drive=MS-D:2, Device=Hard, Bios=0, BootKey=0. These options disable using the function keys to edit booting to the boot menu or to boot to the current in CONFIG.SYS, the "SWITCHES=IF:CN" statement is added. This removes the two second delay after it displays "Starting MS-DOS" and disables using the function keys to do a step by step loading.

Also in CONFIG.SYS is a device named "FOSL.SYS". All this file does is interpret every "ctrl-C" and "ctrl-break" so the user can't boot AUTOEXEC.BAT when AUTOEXEC.BAT loads; it calls a program named FOSA.EXE, which loads FUGGS.386, which is called the "Kuros Guard Corp File System". This is a trick. This is not the file that contains the security layer. I was unable to confirm the claims of FOSL.SYS, who says that FUGGS.386 is a device driver that keeps the Fortres layer on top, not being priority inside Windows. After this is loaded, the classic Fortres boot plays. This is a little more of load, including sounds that play through the PC speaker. This is why if you reboot the machine before properly backing Fortres, everything will know and you will get locked. (This can be turned off by adding "NO" to the FOSA.EXE line in CONFIG.SYS.) If you hold down both shift keys at this time, you will get a password prompt. This will let you disable Fortres for this boot, or put it in diagnostic mode. More on both of these later. Windows then begins to load

and I know for sure the security layer is loaded sometime after the network support is loaded. This is because you can configure Fortres to get its settings for the security layer from a NetWare or NT server. This will put it there I think it loads, and I am fairly certain of my research. KERNEL32.DLL is loaded and then it runs loads and runs MS-GSRS22.EXE, MSOSRV32.EXE, and FORTRES.EXE (the path to FORTRES.EXE was defined in the AUTOEXEC.BAT). This program is called the "Fortres 1.01 Loader" and this is not a file this time. This contains the default file protection settings which can be copied to the settings the FORTRES.EXE loads FORTRES.DLL, which loads the security layer, which is stored in FOCX.WORK.DLL. Note, this file is what we were looking for; the elusive security layer. One of these files, probably FORTRES.EXE, loads the configuration settings from APPMGR.SER, which generates what FOCX.WORK.DLL does. This ends the part that I'm not sure of. After this load is complete, FOCX.EXE is executed and the mouse arrow is moved to the top left corner of the screen. This small alone program simply draws a full window of the FOCX.DWG in the lower right corner over the system tray. Every process started in Windows after MSOSRV32.EXE will have FORTRES.DLL and FOCX.WORK.DLL. This is the basis of my analysis. With these two DLLs, Fortres can screen your actions on every task running. This theory dismisses that of FUGGS.386 being used to monitor all the tasks. Whatsoever theory you want to believe, the truth is every task when MSOSRV32.EXE will have these two DLL files loaded as modules. Anyways, sometimes when the security layer loads, there are Windows loads, EXPLORER.EXE, FOCX.WORK.DLL, etc. This program is the Proxy server for the Base Interface (the part of Fortres (www.fortres.com). This requires the admin to pay for Base as well, and I have never found a computer that is used on. Once this has finished loading, it runs FLOGO.EXE again. The final part of Fortres to load is FOCREPL.EXE, which is executed from the registry in the HKKEY_LOCAL_MACHINE\Software\990-xxxxxx\Windows\CurrentVersion\Run key. Once it is done, FLOGO.EXE runs a third and final time.

Fortres is now loaded and the machine is locked down. By default, Fortres disables the system on all devices so users can't:

- Access Register or Find Files
- Access Shared Areas
- Access My Computer on the desktop
- Access Network Neighborhood
- Access Remote Drives

Access Shared Computer Areas (might check on theory)

Execute Command.com

Start or write EXE, SYS, BAT, PIF, DLL, INF, COM, VXD, OLE, 386, OVL, and LNK files

Format any hardwared

Alter a hard

Make files

Mount or Dis Mount

In addition, Fortres has a list of the programs it will never ever run. By default this list is:

- BACKUP.EXE
- SERVER.EXE
- SETUP.EXE
- INSTALL.EXE
- POLLNET.EXE
- EXPLORER
- DBM06.EXE
- HYPERLINK
- PROGRAMAN.EXE
- MSM01022.EXE
- MSM01027.EXE
- MSM01028.EXE
- PACKAGER.EXE
- DEL787E.EXE
- XCOPY.EXE
- MSCONFIG.EXE
- FCPI
- TSNMGREAB
- MSGETRTS22.EXE

Fortres only checks the name. If I returned REGEDITAEM to HEURREG.EXE, it would run. In addition, Fortres can be set up so there is no saving on a drive or on executing on a drive. If a computer is in a certain mode, then C: only. Less are setting on all drives, and executes from early C. It's basically impossible to do something (they can't warn you). We have already shown that trying to prevent Windows from loading is down head. Over repeating the machine, what's the culture.

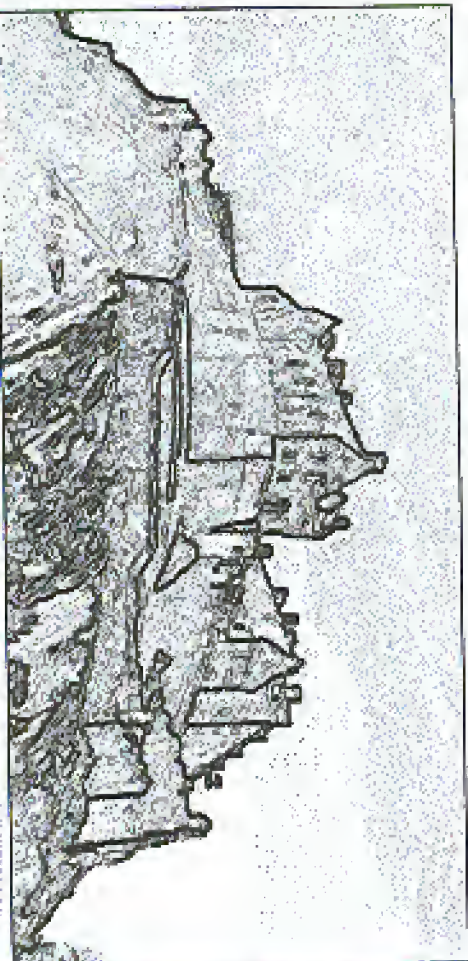
Are you ready for some good news yet? Despite all this security, you can still run most programs. If you go out to the file open level, you can browse the computer. For some reason when the "open file" dialog box is open, you can right click on files and run them. However, programs in the Do Not Run list can't be run this way. Something very useful to you is the, so you can send files off the machine. Plus, even though you can't get to Network Neighborhood, you can still connect to computers on the Network by type "computer name or IP?". However, to do things like security layer specifically process against you need to recognize the Fortres. This is done using APPMGR.EXE. You can run it from Win9x by holding down ctrl+alt+delete, or ctrl+shift+F for NT. Doing this causes a password prompt to pop up. This is exactly

Some words of warning: Fortres 101 logs attempts to access things it controls. These are under the diagnostics window. What the illegal action was and what program tried to do it is recorded. This is meant as a way to see how you need to change Fortres to work with a program. This log is not stored to disk and it does not log the system's reboot. Before you leave, simply clear the log using the clear button in diagnostics just in case. Also in the diagnostics window is the ability to Erland Fortres and you reboot the machine. When this is checked, FORTRES.DLL and LOGNWRK.DLL, which were loaded with every task, are then removed from all tasks and new tasks aren't hindered from. This further supports my theory that though these two DLLs and not FORTRES.SYS, Fortres is able to check everything you are doing. Another thing to fear is an accessory product for Fortres 101 called Citadel Control. It is a remote admin tool that manages several computers joining Fortres 101. It runs on NetWare or NT servers. Currently there is a bug that will not let more than 15 computers be connected to a NetWare server. FORTRES has installed a notice saying they don't know what causes the problem and that they hope to have it fixed by 2000. Needless to say they have not yet fixed it and to my knowledge still discontinue the use of NetWare after version 4. General Control allows the admin to see what each user is doing and issue two types of executors: Policy and Temporal executors. I'm serious - this is what FORTRES calls them. Policy executors include updating the system registry on a machine, starting and stopping tasks, and other admin work. Please note that unlike what APPS/MGR is running locally on the machine, if the configuration is being altered remotely while a person is using the machine, all users' sessions are

terminated. The Temporal executors are things like: erland, shutdown, reboot, and freezing the keyboard. I have also heard an unconfirmed rumor that it can freeze the mouse as well. If any of those things happen to you while you are hacking, walk quickly but calmly away.

The following are things I found that were just weird about Fortres. First, Fortres always runs ETROPTX.E after you close a component of it. This is a great way to make sure a program always runs (only if server unneeded). Also, why does Fortres disable itself when APPS/MGR is used? There must be something APPS/MGR toggles that tells the security layer to take a break. You could easily write your own program that toggles this too. On a different note, FORTRESYS is not a file. Perhaps some thing could rename them before AUTOEXEC.BAT loads? Another little hole is in the FORKINSTN1 file. To make installation faster, the file allows for several logs, one of which is "password...". Who knows, maybe someone was stupid enough to use it. With Win NT, simply logging in as "Administrator" will giveable features not yet log out. Finally, when I view in the more likely exploit is MISC/SK32.EXE. This file seems to be inside the security layer since it is loaded after the kernel and itself loads Fortres. Perhaps it could be used to create a link before the security layer, and thus allow you to do what you want. (Again, I do not know what that is password protected might be a good thing here.)

I hope you better understand Fortres. It really is a well written program. Any system admins out there who want help on how they can configure Fortres 101 on their machines are welcome to e-mail me and I will gladly help them. Rock on.



Passport Hacking

by CHRIS SHIMMEL
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This article introduces a security vulnerability in Microsoft Passport. Specific details explaining how to compromise your Passport account as well as counter-measures to do this will be given. However, this information is intended to be used as a technical example. The objective is to give a detailed analysis of security on the web while illustrating some common misconceptions. I combine with some suggestions for using the existing Passport mechanism as well as ways to improve its security.

Background

Microsoft Passport is a mechanism designed to allow users easier access to services offered over the Internet than a unique registration. The intent is that users may register for a Passport and then use services on various sites without having to register on each individual site. Search is a hassle for the user in terms of time spent as well as continued password maintenance.

An attractive feature of Passport is the WebSite Helper. WebSite Helper allows you to save credit card information to add to the personal information you may collect. This can be useful in participating sites to make purchases. Future references to the Passport mechanism apply to both the helper itself and the WebSite Helper.

Introduction

A Passport is merely a collection of cookies saved on a user's computer. These cookies identify the user on a Passport enabled site. There is no server to server communication involved in the Passport mechanism; all communication is handled through the user.

The various cookies are described in the following table:

name	description	secure?	path	storage
BrowserTest	passport.com	No	/	memory
MSPVIE	passport.com	No	/	disk
MSPIDom	passport.com	No	/	disk
MSPAuth	passport.com	No	/	disk
MSPProd	passport.com	No	/	disk
MSPSPe	passport.com	Yes	/password	disk
MSPProq	login.passport.com	No	/	memory
MSPVIE	webtel.passport.com	Yes	/	memory
MSPSTok	webtel.passport.com	Yes	/	memory

Microsoft Internet Explorer versions prior to 5.5 have a cookie vulnerability that allows client side scripts to reveal information stored in cookies not intended to be shared with the current web server. Since the mechanism behind the helper is based entirely on cookies, the problems of this combination are obvious. The most troubling result of my research has been the lack of major obstacles to complete impersonation.

Security Test Score

A check it only as strong as its weakest link. Sound familiar? It is astounding how many people are given a false sense of security because something is encrypted. The majority of the cookies mentioned in the previous section have cryptical values. You will notice that no attempt is made to describe these values in their entirety. Why not? Well, frankly, because it is difficult to do and absolutely unnecessary. While cryptographers can be a very challenging academic subject to pursue, the point of this article is to show how easily a mechanism such as Passport can be cracked. Someone attempting to crack into a web site's web service will generally take the easiest route possible. This should be the reason you care.

When a user goes to a Passport enabled site, the site itself does the decryption. By simply presenting the same encrypted variables to the site as the legitimate user, anyone can impersonate that user. The only reason anyone would need to be able to decrypt the values in the cookies would be to create a Passport enabled site.

Now, if this all sounds familiar/eggy/familiar, it is because it addresses the same absurd misconceptions that have brought about the RSA/SSA theorem. Re-remembering the contents of a PVD is only necessary if you want to play the DVD. For those wanting to copy the DVD, downloading does absolutely nothing to help. The fact that some people do not understand that is not nearly as shocking as the attempts of the RSA/SSA (and its supporters) to lower this ignorance.

HTTP

Most people who browse the web are familiar with HTTP, though the details are not often taught or understood in the case of many web developers. To understand how two sites communicate, Microsoft Passport, a basic understanding of HTTP 1.1 and how it handles cookies is required. There are many details of HTTP that this article will not discuss. Please refer to RFC 2615 for the full specification.

The basic HTTP scenario is a simple transaction consisting of a request and a response. When a user is browsing the web, the web browser (client) makes a series of requests to the various web servers around the Internet that the user vis-

For those who might have gone all out and purchased a "Plus" receiver, you will want to edit starting at address 83C8 and enter: 05 3x 3x 3x 3x 20 207. You will want to put your zip code where the x's are. So for a New York City zip, you would put 55 31 50 30 30 31 50 20 or zip code 10001.

Once completed, save the eeprom file as H04011A. Remember, this needs to be from a valid H card since the emulator will verify from this binary. The H card inserted into the programmer does not need to be valid from base or not, since its only purpose is for decryption.

Before we ALEX the H card, we will want to use basic H to "one step clean" the card to 28 updates. Once this is done, you are done with Hatched until you decided to further experiment.

Now, on to creating the ALEX seed. This tutorial continues your H card in and as a go-between in the emulation process. You will need to install WinDySlex.exe first and then you will be able to open the TurboALEX.vch script. (Use the Source Table) on this one. The .vch file has instructions and information contained at the beginning. Using the program is not difficult, but sometimes being informed is so useful.

Go ahead and execute the TurboALEX.vch script using WinDySlex. A small window will open up. For those Terrik fans you will not be disappointed. For the rest of you, expect to be disappointed with the GUI. From the new window you will want to click the ALEX button. The ALEXing process is the most time sensitive of them all. So be sure you have nothing running in the background if you have a slower system.

Finally, you will need to create a DOS boot disk. Under Windows 95/98 it is easy to do from the control panel. You are on your own here though. Once the disk is created, you will want to copy the modified binary h2600.bin to this disk. Following that, copy the emulator software to the disk from the SLH41 archive.

You need to copy the s1644.exe file to the boot floppy. You will also want to read the readme301.txt file in the zip archive to see what other command line switches there are and some things you can do to troubleshoot timing issues.

Now boot your computer from the boot floppy. You can automate the next steps with an autoboot.bat if you are so inclined. For this example, we will assume your HSD card is plugged into slot and the ISO programmer is

connected to COM2. Your boot floppy should contain system files for boot, the h2600.bin binary, and the s1644.exe emulator application. Thanks to Pierre G. Marinneau (HGM), the dry week of the communication between programmer, computer, and receiver is squared away. You will want to power off your DTV receiver at this point. Make sure the ALEXed H card is in the programmer. Once you are booted,

type:
A>cs abcdefg h i j k l m n o p q r s t u v w x y z
Give it a moment, and then power on your receiver and change to channel 100. Give it a little bit so that the emulator can sync with channel 100. I can not say if spending with 100 is necessary. There are mixed reviews on this and there is some information saying that 100 sends initial seed information. True or not, I have found that channel 100 has fixed things in the past. So it is recommended. Also, since HSD card fragging might be required. You're a hacker, figure it out.

If things are working, because of the 6's command line switch, you will see the communication between the computer and the receiver. Spend some time and watch what is going on. If you check the forums on aldris, hark or some of the other online forums, you will find commentary about some of the more interesting streams.

These forums usually discuss ThreeJ's attempts to kill modified H and H0 cards. In this case, you should be just fine. Because the emulator maps the signals being written to the card and just updates them in the h2600.bin binary in memory, your H card is safe. You will want to type "q" in order to quit the s1644.exe application so that updates are saved to the disk for the next time you boot.

Some Notes
From the emulating: "Thus far, the only universally emulator-compatible HSD known is the Hughes H1 series. However, some vendors have reportedly not been able to work with Hughes B2 series HSDs and RCAAD22 series HSDs."

For those looking for something that will run under Linux, look for a file named p16u-0101-0011001.targ on the net. If you have trouble finding this on the web, you might want to read to the end of this article.

The H0 card is the next version of the H card. Currently emulation is not possible for H0 cards. These cards are susceptible to the ECM, sent from on high.

Additional Resources

At the time of writing this article, a post to slashdot.org about a version of the emulations software for Linux allowing for distributed network sharing of a single H card by many receivers brought a flood of legal attention to the DSS enthusiast community. The application of law is called *Dieci* written by mrg343. He has documented his work on the protocol line as legal threats under the DMCA. Also, the moderator because of threats of litigation, took one of the best resources for the DSS enthusiast (www.backto.com) down. Oddly, this site is hosted in Canada by a Canadian but likely due to NAFTA trade treaty status, the specifier of legal threats from the United States is able to affect the Canadian citizen.

In addition, many of the text files and binaries mentioned in this article are becoming harder to find. As consumers, there are reasons

for fair use of this technology. For the terminally curious, access to how this equipment works is invaluable to the psyche. I myself recommend anyone playing with this technology get a subscription to *DieciTV*, if not the most basic package. We as hackers are not here to cheat. We are curious and our desire to investigate and discuss this curiously should not be a crime.

Anyhow, you can still find more resources on this topic at www.usanodprogram.com and www.dsslist.com. One of the best resources I found while constructing this document was p16u-research.zip. This file has a hardcopy of articles, test files, tech sheets, and other documentation that mrg343 used to develop his application. People have made CDs to post all of those files and information on the distributed pip not made possible by Grunella.

Enjoy your television and try to learn something about television by hacking DSS.

CodeRed 2

op how to anonymously get root on 250,000 machines overnight

by Braddock Csekell
braddock@braddock.com

This article describes a mess through which a complete list of the estimated 250,000 CodeRed II infected and backbone compromised hosts can be easily obtained by any individual who has been keeping a web server log of attempts on his machine, by using the backdoors on the machines that have attacked him to obtain the web logs of the infected attacking HIS web servers to learn of new infected hosts.

The strong recommendation from this report is that as part of any CodeRed II recovery effort, the system web logs should immediately be destroyed, and Intrusion Detection Systems should be checking for and tracing recursive attempts to access web logs through the backbone.

The CodeRed II worm has been infecting HIS web servers with a speed equal to or greater than that of the original CodeRed. The original CodeRed infected what is thought to be all vulnerable machines, approximately 250,000 hosts, in under 24 hours.

While CodeRed I was relatively harmless,

CodeRed II installs a full administrator-access back door shell that can be accessed via http. This creates a very interesting situation and, with the techniques discussed in this article, opens a new potential door for mass system cracking.

The problem with releasing a worm or virus to obtain some information of value is that to transmit the information back to the worm originator a very clear trail is created that can be traced back to the perpetrator. Primitive and naive worms or viruses sometimes attempt to e-mail or otherwise communicate password files or information back to some originator point, allowing a trace to the original author. A more sophisticated worm might attempt to just pass information upstream to get it closer to some originator node, and make attempts to destroy records of the transmission. But this too leaves a trace of the worm's spread. All records of the transmission in things like firewall logs and IDS systems can never be removed.

It is difficult enough to find an anonymous enough node to make the initial release of the

worm. Presumably one would do this far from home in a previously unpatronized Internet cafe or the like, through a large number of randomly cracked systems. If an author actually uploads some attempt to "return to the scene of the crime" to explore anything of value, the worm might send back to some anonymous node. It would most certainly be caught.

The alternative to this is to attempt to make the information the victim gathers public, and then attempt to retrieve it just like thousands of others will. For example, a search might send password lists to a Usenet newsgroup or post it in some public forum. But any public forum usually has some form of moderation and administration, so any malicious information at such a site would not stay online for long.

In addition, the more sophisticated the initial worm, the more stylistic and linguistic "fingerprints" the original author will leave on it. Posting to public forums may well double the code in a single worm. If an author has ever made any of this code public, there may well be governmental agencies that could use code fingerprinting to narrow the field of suspects, particularly if other profiling information can be used.

If a more "anonymous connection carrier" system like FaceNet is ever successfully put in place, this may well change the landscape. But true anonymity will probably always remain elusive concerning national security or currency laundering enforceability. It is at stake, even if enforceable. Deasonian legal means are required to achieve it.

CodeRed II, however, presents a very different alternative. CR2 infects its hosts with a simple worm, inserts a simple administrator-access backdoor shell into the victim, and begins scanning for new victims. At first glance, the backdoor is of little use to the worm originator. After all, the originator has no list of infected hosts communicated back to him or left at some server drop point. The originator, like anyone else, can perform massive network scans for the backdoor, but that would put him on a relatively short and easily compiled list of suspects. The worm also keeps no log of hosts that it has infected, and indeed no log is essential to keep the spread untraceable to the originating node. Perhaps a public key encrypted log could be compiled, but that leaves us back to the original problem of a fixed "drop point" or compromise of the data.

Lack of usefulness appears to be the case.

except for the fact that the Internet is now saturated with CR2 worms, each leaving web logs across the Internet full of records of buffer overflow exploits, with the infected host's IP address. These stack attempts point to an address that requires a massive Internet search; they serve to announce the infection of the attacking host. And they do so in a way that leaves no direct trail of initial spread of the worm, and thus no direct risk of discovering the originating node.

This means that by the end of the first week, I personally had in my web log the IP addresses of over 100 random hosts with full-access backdoors installed that I could attack directly. One hundred hosts on different unrelated networks is a large compromise, but not something that requires a massive Internet search to achieve. This is not enough to make the plague of a worm worthwhile to its originator.

However, each of those 100 random infected hosts I know about are also 100 web servers with logs of, for example, another 100 random infected hosts each that attempted to re-infect them. That means by breaching into the 100 hosts I know about and reading their logs, I now have backdoor access to approximately 100x100=10,000 hosts. Repeat this another level (probably obfuscating from the books nodes), and I will have 1,000,000 break-in attempts by random hosts. At this point many of these attempts will be from duplicate hosts, since only an estimated 250,000 hosts will be infected (this from the CR1 estimates), however it is clear that the implication of this worm is far greater than random hosts with backdoors. It provides a clear mechanism for obtaining a list of thousands of infected hosts with backdoors.

While this technique is nice, it is still not entirely untraceable. IDS systems will surely be looking for this type of backdoor exploiting traffic in the usual form, and contacting several thousand hosts either directly or through a worm backdoor distributed mechanism will be detectable on some level. A full list would require the recursive retrieval of web logs from several thousand hosts. However, the originator of the worm himself does not need to fear exposure... he has essentially made this information available to anyone who understands CodeRed II and its implications described above. A public list of all infected hosts is probably already available online.

REPUBLICAN THE RULES - The Ultimate DRM Hack

by The Matrix

"It's like to be verifiably integrated from the moment of creation right through to the moment of delivery" - Rupert Murdoch

Sherry said to us: "We've been spending the very thing the most powerful media empire on the planet has ever with all his heart. While our intrusive hacker live buzzes with wireless activities ranging from encoding DVD encryption to co-opting Microsoft's Media Encoder into Divx, the real power to publish is being systematically and subversively removed from our economic grasp. The power to seal is overwhelming the power to share, and the real winners will be our children."

What am I haranguing you about? Digital Rights Management, simultaneously the most liberating and oppressive concept within the modern computer world. DRM, as it's usually known, is a class of computer systems that control access to data in its digital form. These systems are copyright, expensive, and represent Big Media's best sword against the notion of a fair, non-profit oriented means to get creative data to its consumer. Just the company names in this space sound ominous: Prometheus, Lockstream, Microsoft.

So what, you say - we'll hack the things. Rupert, Jensen, Gates... they don't stand a chance. Well, maybe we will and maybe we won't, but the point here is not about hacking your way to free copies of Star Trek - The Next Generation. It's about telling the story of how you did it. And about getting your story published and distributed the way you want it distributed, not the way Big Media wants to do it.

Let's get real here - any reasonably sophisticated DRM has a few common components:

1. A device registry: These exist because Big Media wants to control where the creative data actually goes and how long it stays there. (Just with me, no problem, so long as it's their creative data. But what about my creative data? What's going to control where that goes? Let's

unhooked, the answer is Big Media or nobody. Nobody may sound like a solemn answer to you, but if you spend six years creating this data, you may want to at least get acknowledgment from people who are using it and like it. And maybe you'd rather not bother them after they've acquired it onto their own device. Plus, maybe you included some sort of value creation (like an advertisement or subscription) that is assessed based on the number of consumers you have. Maybe you buy lead from the providers.

2. Encryption. Generally, the gather-streets stuff, as it's reasonably difficult to hack. It's a waste of time to hack it anyway, as the back door is usually left wide open during the events that occur on a computer after stuff has been decrypted. It's actually this little back door problem that is at the root of the most oppressive aspects of DRM: It leads to the design and construction of electronic devices that embed a private enterprise's approach to controlling any data that shows up on that device. Again, cool with us - it's their device and if I don't want it, don't buy it. But what about my creative data?

3. A Package: A package takes the creative data and packages it for distribution under DRM control. They often embed some little features like the ability to create a same-name program that on a target computer, can access Operating System necessary space and perform intrusive, privileged acts like deleting data. Again, cool with me so long as it's their data and I generated them, through some sort of license, the right to do so. But, do you think for a second that it's some citizen such as yourself could afford to use such a powerful tool? Think again.

4. Keys: Every DRM has keys. Keys are often hackable, but they are also immensely powerful mechanisms for enabling a prescribed sequence of events to occur on millions of computers. Why would we even consider having such power in the hands of private enterprises? DRM companies are undergoing their first round of shakeups, and as any futurists (101

Student knows, only a few will be left standing. DRM is a commodity, which means - under the track currently underway - one company will eventually dominate (think Microsoft). The notion that consumers will use multiple DRMs based on which creative data they choose to purchase is ludicrous. The architectural underpinnings of these systems are just too weak, which translates into too many bugs and too many hoops for the Average Joe to jump through to use the creative data. Heard any BlackMatter music tracks lately? I didn't think so - and neither does BlackMatter.

There's a massive issue at stake here - the opportunity (just the right) for individuals to obtain the snore, or better, level of DRM capability as the big boys and, in the process, to make sure the one DRM left standing is as robust as possible and provides equal opportunity for all. Just like Linux, FTP, or T-shirt.

decoding COPY PROTECTION

by Park Rat Sam

I am not a music pirate. I am a Canadian though. Oh! Man! I, so all those nasty DMCA rules don't (currently) apply to me.

I own a few licensed CDs, all 100% percent legally ripped and encoded. I don't even own a system to even just my computer. (Kidding my computer up to date enough to play the latest games is plenty expensive enough as it is!) When I buy a new CD, the first thing I do when I get it home is drop it into my computer, rip all the tracks, and encode to MP3.

Imagine my horror when I bought a disc that said:

"This audio CD is protected by SmartCopy (TM) MediaClon (TM) Ver 1.0. It is designed to play in standard Audio CD players only and is not intended for use in DVD players."

And sure enough, if I put the disc into my drive, all my ripping program could see was a bunch of data tracks. I had to beg, plead, and borrow to use somebody else's portable CD player just to hear to one file.

How MediaClon Works

MediaClon supposedly "protects" audio CDs in two ways:

1) Deletable errors are introduced into the audio stream so that ripping programs introduce gaps and clicks into the ripped data. Normal CD players have circuitry designed to cope (erase or best) with corrupt data, such as caused by scratched discs, so they can interpret the missing data with the best error rate possible.

2) The tracks are marked as data tracks so that a computer won't recognize them as audio. All of the audio data is still there, laid out on the disc exactly as you would expect, except that you can't pick a track and select "Play". Somehow this didn't seem like me at least say more of a protection than that lame "Copy Protected" bit that all CD ripper already ignore. If you could just point your ripper at all the right sectors, all the audio data is sitting pretty right there for you. Obviously marked and unscrapped.



It's time for a call to arms. It's time to petition the IETF to develop an open protocol for the common elements of DRM. It's time to distinguish the common elements from the value added elements and to create a framework for the competitive circus that now exists in the DRM marketplace. It's time to donate our skills and abilities towards the creation of this system and to use our hacking skills to break it and to fix it. It's time to wrestle the power to publish and control distribution of creative data away from the hands of a few individuals and into the hands of the Internet user. It's time to educate our children that the opportunity to publish and compete with Big Media is theirs and the right to consume is limited by ethical behavior. Soon, it will be too late.

The technology is close enough; it's now about economics, sociology, and setting objectives. Make your opinion heard.

This cannot qualify as a "protection" device, because CD-ROM discs are designed to read raw sectors from a disc. The only thing preventing ripping programs from obtaining information from data tracks is that the data stored in non-audio tracks is not actually audio, and you wouldn't want to use those bits as your speakers by ripping random noise through them!

cdparanoia

My favorite CD ripper under Linux has got to be cdparanoia, because under the GPL, it's huge, gives us a whole program. This program was already designed to deal well with scratched CDs, so protection marked one above is already dealt with effectively. The only thing left to work around is the data tracks which have a few code snippets from cdparanoia, with the offending lines marked:

```
cdparanoia-III-alpha&Alpha.h:
line: code
```

```
999 : switch(cddg_&enddij)
```

```
908 : default;
```

```
909 : report("unlabeled to open disc");
```

```
910 : exit(0);
```

```
911 | }
```

Function calls report returns "-409" if it cannot locate any audio tracks on the disc, which causes cdparanoia to die here.

line: code

```
1000 : front-track110e-tracek21h | }
```

```
1001 : idf(sddg_track_audioptr000)
```

```
1012 : report("Selected span contains non audio
```

```
tracks. Abort(sg_start);
```

```
* 1013 : exit(0);
```

```
1014 : }
```

This section of the code is similar except that here it is verifying, one by one, that each of the tracks you're trying to rip is marked as audio. Any data tracks in the bunch and cdparanoia dies. Bypass both of the lines marked ***, and cdparanoia will happily read any sector on any disc, whether marked as data or audio. The patch below adds a command-line option "-N*" to do just that.

```
--- main.c.orig Sat Aug 11 16:52:25 2001
+++ main.c Sat Aug 11 16:52:31 2001
@@ -586,5 +586,5 @@
 |
```

```
-const char *qstring = "eoc:nc:hd:dg:s:gr:Rn:atv:Y:Qb:Zz:z:YXWBI:TE?";
+const char *qstring = "eoc:nc:hd:dg:s:gr:Rn:atv:Y:Qb:Zz:z:YXWBI:TE?";
```

```
struct option options [] = {
@ @ -662,4 +662,5 @ @
int qopt, only=0;
int batch=0;
+ int Store(CdQ=0);
```

```
/* full paranoia, but allow skipping */
```

```
@ @ -791,4 +792,7 @ @
sample_offset=cdinfo(parg);
tracks;
```




```

+ case 'M':
+   Modem(CIO-1);
+   break;
+   default:
+     usage(stderr);
+ @ -1906.4 +1907.7 @@
+ case 0:
+   break;
+ case -403:
+   if Modem(100);
+   break;
+ default:
+   report "unhandle to open disc %i"
+   out 1;
+
+ J0134yda(CIO0)
+ fort1=cach1jpe119e8241++1
+ file=dda_cwork_sool1opt1j0f1

```



TIME TRAVEL

by **It's AAD!**

itsasabi@yahoo.com

Perhaps you've scanned the telephone numbers in your area, so now comes the time to alert concerning the unknown. Each time you found a number, number your software should have logged it for a later attempt.

Some numbers will be fax machines - not a lot of fun there unless you have some product you may be selling. Other numbers may be sales (insurance, 162). Perhaps you may discover an X10 house system.

But not so well known are timeclocks. And if you find one you may want to know how to get connected and enter with administrator privileges. That is exactly what you may learn from reading this. But first the disclaimer: Please don't try this at home kids, as you may damage the stored data held in the memory of the timeclock and maybe someone will not get paid with that out of the way, or to a brief description of a release (backlog) (see Phase 1). Mounted on the wall, employees "swipe" their



cards, much like a debit or credit card machine. The reader then decodes the simple Mhz-classic encoded magnetic strip. This information is typically the employee card number. The timeclock then saves this number along with the date, hour, minute, and seconds for later retrieval.

The timeclocks I am discussing are "TTC" made by opex systems in Utah. They currently

make several models, including "Bioromatic." The one in Phase 1 is a model "94100," one of the least expensive units available. It can be remotely accessed via a 2400 baud modem and typically will have a dedicated land line.

Normally payroll departments will use the provided Windows program to download the punches from the timeclocks every two weeks. They should be the only ones who call the timeclocks. However, using a simple terminal scripting program like ProComm, anyone can access the timeclocks.

To enter into communications mode with the timeclocks you will need to set your speed to 2400. Use N81 on the M1090s. The 10250 (see Phase 2) require 55.5333 baud also N81. Then, when prompted, enter the username (default is the clock number). Next you will be asked for the password. Enter H1C (usually, the password is seldom changed. Another way to access these clocks is through direct "rally chair" as the timeclocks all come equipped with an RS232C socket in the form of an RJ11 jack next to the one marked "tel." For this the company provides an adapter to your serial port.

The 10250 also has a keypad. When you press the arrow key you can gain "supervisor" or "user" mode. The default password for user mode is 22222. This will allow you to read all the card numbers and punches. If you enter 11111 you will gain supervisor access and you can change the time, or even clear the memory of the punches.

Myths about TCP spoofing

by **Grandmaster Plague**

To many IS&T hackers and aspiring hackers, the myth is perpetuated that the surest way to not get caught at whatever it is you're doing on the Internet is to "spoof" your IP address. Fully intended to clarify this belief and give examples of when spoofing can be best used.

What Is It?

"Spoofing" is a process by which the IP address of your machine is made to appear different from what it really is. The purpose of this is

to access the clock with a source that will cause the modem to hang, meaning it will not reset. And after many calls with tech support, the only remedy is to disconnect the power for 15 seconds. This can cause a big headache for the company that is trying to track their employees' time. Other versions don't seem to suffer this "hangover."

I am sure there are other links on the market and they no doubt will use the same type of system, typically a programmable CPU like a Zyneg or Intel 8531, with a simple modem and "buggy" software.

As a note if never ceases to amaze me that those companies all like to use simple passwords, like their company name, store number, etc. How long will it take before real security is the norm? Only time will tell.



so as to hide your true point of origin. Examples: if your real IP is 138.13.235.182, and you spoofed it to 199.199.199.199, then your IP address would show up as 199.199.199.199 in the remote machine's logs. Thus your real IP address would be unknown. Many novices (and others) think that if they get a magical "IP Spoofing" program which modifies the source IP address that maybe the source MAC address field of each outgoing packet, that nobody on the Internet will know what their true IP address is.

But Wait....

The problem with this belief is that TCP (and most other network protocols) is a two-way street. This means that for just about everything you send out to a computer on a network, you expect a response back. This is a problem because if the remote machine thinks that your IP address is 199.199.199.199 and your address is really not, then the machine will try to send information back to that spoofed address and you won't get the information (because it's not your address).

TCP Specifics

If you still think that you can use IP Spoofing for the "one-way" protocols (like rsh, rcp, etc.) on the Internet, think again. The problem is that if you want to be connected to the Internet, your machine must speak TCP/IP. TCP/IP is the foundation for the Internet, thus, every higher level protocol (such as HTTP, FTP, etc.) must use TCP/IP. TCP/IP gets information from point A to point B. What happens when it gets there is the responsibility of higher level protocols. Now the reason that this is a problem is that TCP has a built-in "feature" that makes sure information is going to and from the right place. This is called the "TCP three-way handshake." Basically, it makes every Internet communication a two-way street. Here's how it works: Assume machine A and machine B are starting a communication. Machine A says, "I'm machine A," machine B responds, "I'm machine B, you say you're machine A?" Machine A then responds, "Yes machine B, I'm machine A." A packet must pass this little test in order to be received by machine B. As you can see, all communication on the Internet gets turned into a two-way street.

Two Solutions

There are two simple solutions to this. The first solution is a form of one-way communication called "Blind Spoofing." The theory behind blind spoofing leads down to timing. Essentially, a machine that's called XYZZ fakes the TCP three-way handshake by saying, "I'm machine FOO," then waiting for a bit as machine B responds to the real machine FOO, then saying, "Yes machine B, I'm machine FOO." The real machine FOO won't know what's going on because it will just ignore the packets that machine B sends to it, thinking that machine B is its name. Machine B won't know

what's going on because it's receiving responses from machine FOO (which are really coming from machine XYZZ). So machine XYZZ has fooled machine B into thinking that it is really machine FOO and it thus passes the three-way handshake. This can only work well in one-way settings where it is not necessary that the client get any feedback from the server. An example of this is SMTP. You could blindly spoof your IP address to a SMTP server (to make it think that you're an Internet IP), and thus get your mail message sent to someone else with a different originating IP.

The second solution to this is a little bit tricky. It's the best way to spoof when you want information back from the server. This solution is called "Active Spoofing." Active spoofing leads down to blind spoofing, but at the same time you are sniffing communications going back to the spoofed host. Using the example above, you are also sniffing the packets going from machine B to machine FOO. In order for this to work, you must either be on the same hubbed subnet as machine FOO or you can do some route table modification to get the information to pass through your machine. You then watch what machine B sends to machine FOO for the entire session. This is an extremely complicated process and changes from protocol to protocol. Currently, I am not aware of any tools that automate this process.

Conclusion

Spoofing isn't really all that stacked up to be. It isn't the be all to end all of covering your tracks. It does have its interesting uses (including false mail, reverse, and more), but is extremely difficult to implement if you want information back from the target host. If you really want to cover your tracks, it's better to reuse all your traffic through some wigwags (or something). There are loads of IP spoofers out there. Some are more useful than others. If you want to track up your own spoofer you can use rawsocks. Alternatively, you can use spoof (a spoofing library) available at: <http://hacking.linux.no/code/patch-tools/index.html>. For more information on spoofing, read *Howe's Penetration How Network available from Syngress Books* (Chapter 11 is all about spoofing).

By Gary (Gary)

Playing with Qwest DSL

By Phobos

I was at a friend's house and he introduced me to an IP range that got my attention. Every address within 63.224.0.x that we calculated to give us the same "chess" prompt. My sense of curiosity immediately kicked in. So what did I find? Based on...

As it turns out, the routers were manufactured by Cisco and are a part of Qwest's DSL network. Each router is placed in a customer's home and quite curiously configured to their type of service. These routers have three basic access levels: enable, and debug. When you initially log on to the router, it asks for a password. On 80 percent of these things, there isn't one. So just hit enter and you're into enable mode.

So what can we amuse ourselves with from here? Besides passwords and ping utilities, there is a reboot command, and a couple of commands for getting info on the router's configuration. Typing show of those ("show" or "show") gives you a list of arguments to append with definitions. Pretty simple, eh?

Once we got tired checking out the configuration info, we can move up an access level by typing "enable". Again, Qwest is lazy and has neglected to set up passwords on the vast majority of these routers. Besides all the exec level commands, we now have access to "set" and "write". Briefly, "set" allows you to change the router's configuration file, and "write" writes the file to the router's NVRAM. After the file is written to the NVRAM, you must reboot the router to activate any changes. One thing you may want to check out before you make any and fashions here is if the system is active ("show system"). If it is, everything you do is being logged to a remote system. Just disable it using "set system disabled". I've never seen anyone access a router with this feature enabled, but it's worth looking on.

One interesting thing that you may have noticed you can't change the router's IP address within enable. This is done through the following command: "set interface wand ip [ipaddress] [netmask]".

In place of the first set of brackets, you choose to change either the upstream or downstream to whatever value you entered in place of the second set of brackets. The band pass will automatically adjust to match your settings. You may also want to play with the router's transmitting power with "set interface wand tpower" or "set interface wand rpower".

What's I won't go into much detail about many of the specific commands (the router's software makes it easy to figure out), there is one more thing I would like to point out. The software fringe and debug console file can be downloaded, changed, and uploaded again using the "TFTP" protocol. All you need to do is make sure TFTP is enabled ("set ftp enabled") and then you can connect using a TFTP client. For those of you unfamiliar with TFTP, there are net directory services, so you must know the exact filename of the file. Lucky for you, I've already done the research. The software image is either named 66756.x.x.bin or 66766.x.x where x.x.x is the version number. As of the writing, 2.2.20 was the most current version. The config file is stored as config.cfg. Just remember that any modifications you make will not become visible until you see the error from manual exec reboot.

Lastly, I'd point out the debug mode. It has all kinds of nifty commands for testing the Qwest network. I'm not going to go into detail about any of it because if you do know it, you're smart enough to have a clue what you're doing in there. You don't need any of my help. I'm sure there's some interesting things hidden in there, though.

That's basically it. If you'd just read that I do any change liability disclaimer. I contacted the company about this problem way back when the routers were still 185West property and nothing was done about it. So maybe this will gather them to wake up. A good resource for learning more about these things is the Cisco website. Just search for either "CBOS" or "Cisco 636 router". Good luck!

(How? TV 2.7.11 is my exact, and accurate information, along with many other names. For network information you can view what digital hub you are connected to in your neighborhood, your cable box's IP address, name, paired IP, and a little more. With this information revealed the possibilities are endless for you hackers. It even gives you the option to name your channel options - if you have parental control disabled - as if the names are supposed to be changing then cable box and accessing this menu, it allows you to change what frequency your box is tuned to - why I don't know. That is the best thing the cable company would want you to do, right? However, physical around much with digital cable, but before digital you could reach PTV channels and premium channels with a frequency modulation on your cable wire. Interesting what possibilities are present with me as they to change the frequency, right in your box. If you turn them into when in this menu, it says 6111 on the screen. Now if you have this made by handing off your cable box, trying to simply enter 6111 will not work. Now if the cable company wanted you to access this menu, by asking to disable parental control, they wouldn't have tried to hide this diagnosis and system information element. It says: finds our master about anything I would like to know.

phly

Dear 2600:

After reading Ed's letter in 18.2, I had to share in. While in college I worked for Radio Shack. I was fired for allowing my phone number percentage (the amount of phone numbers required for maintaining a position at the Radio Shack) to fall below 92 percent. Since I only worked on the weekend, if more than one customer asked me not to handle their information, that I immediately left below the quota. I suppose I could have found some workarounds, but I don't remember how I hacked a six figure position in the tech industry and then share and well in LA.

As a side note, each Radio Shack collects customer information on a daily basis and uploads it to a secure server at the home office in Texas at the end of each night. This is done via a dialup 56K 135K point to point connection when the store manager closes out for the evening. As part of the process, the manager is given a journal which includes the activity for that transaction which is typically filed away for archiving. As time goes by these journals become cumbersome and are supposed to be shredded. In my experience through the four stores that I had worked at, this shredding typically just drove them out.

This shredding data is used for a myriad of things. The average store always held on that when confronted by customers as to why we ask for this we want to "just tell them that it's collected for this we save your from a archive." We know that this is an unsecured server since you have to either make a purchase to get a "free" catalog or make a formal request. I can only imagine how valuable this user data would be to other vendors. Ever notice the arrival of a Cloroxed catalog after making a purchase at Radio Shack?

Needless to say, they take collecting this data very seriously. The next time you go in to buy a capacitor

and are galled for not returning on where you live and what your favorite color is, refuse and watch the clerk's temperature rise. There is an actual pressure put on these folks to gather data. But they will often add fake customer info to your receipt if you decline so that they can come back to work the next day. In practice that I refused to do and thus was released from employment. What do you expect from a company that pays its employees \$4.25 an hour?

hacker181

Dear 2600:

I've never had great luck finding more than one way to expose the term "hacker." Led to some success reading a new word compiler my eye. I figured looking up the term in a dictionary. I gathered the impression that a dictionary was a private web page very hard to edit with technology and computers in general. I don't know if anyone has expressed this love before, but I was very pleased to find another possible way to expose "hacker."

Don Sherris

Dear 2600:

I was just sitting on the computer flipping through 18.2, and I saw a letter from "Anonymous" (name that guy wishes a ton of stuff all over the place, it comes from a previous article about getting off a telemarketer's call list).

He says that the Do Not Call calls happen to a "telemarketer" whether it's a surveyor, salesperson, or a fund-raiser. This is in fact untrue. The laws are actually very specifically geared towards calls of a sales solicitation nature, such as calls trying to get you to buy a product, be it a new TV, vacation homes, magazines, etc. This means that "surveyors" they gather the term "telemarketer" so well as fund raisers are exempt from the laws.

It may also have been incorrect in saying that the Do Not Call lists are company based and not offer based. This aspect varies from state to state, so you can't lay down a blanket statement. The existing company is also only directly responsible for applying the laws to the states in which they conduct business. This is another loophole as states define "existing business" differently. Some states consider it only where you have a physical presence (i.e., where is the telemarketer selling or a phone), other states consider it anywhere that the calls go to (i.e., where is the person answering the phone call).

There is not specific wording needed to be placed on a "Do Not Call" list. As long as you make it clear that you do not want to be called ever again. "Please no you do not want the please" as well as "Yes, I wish, don't ever call me again or I will kill you" can still hash surface (although the latter may get you in trouble because you placed off a person who may have lost of personal info about you sitting in front of them).

In the event of a company wide ban, the company is responsible for making sure the phone number in question is never called again for any jobs. If that means retying it from this, fine. If it means setting up a proactive dialer to bump the missing numbers,

that's good too. Whatever method they want to use as long as the phone number in question is never called again. Notice I specifically chose number. If you have more than one number, all bets are off. They can call sub and on any one of your numbers, and you will need to inform them for each number individually.

For sub-based dialing, they just need to remove you from the one offer in question. Future offers are just fine to call you on. Job-based dialing is the name customer of the two in the laws I have seen.

I might, however, argue what is interested in dialing with this to read up on 1800 state's laws. Many telemarketers don't bother for don't have the power to record a number as to do not call. So if they call you back, you might be able to collect fines. Many states offer reparation in the range of \$100-\$500 per call in violation. It is up to the disposition of the calls to show proof that they are agreed to not be called, and show proof that they in fact were called again in a manner that violates the law. Tape recorders work well for this, but again, ask a local lawyer. Not all states allow you to record a conversation without consent from all involved parties.

Just some thoughts from someone who hasn't spent a summer or even "over a year" working in telemarketing, but rather has spent the last 20 some odd years of my life dealing with the technical and administrative aspects of setting up and running call centers throughout the United States.

Dear 2600:

In issue 18.2, Mike G. asks where the physical files can be found now that they are no longer maintained at pen.com. You can find the entire archive at www.pen.com. They are currently looking for a new business. Any volunteers?

Miguel

Dear 2600:

The response to Jeff's letter in issue 18.2 about giving Radio Shack's corporate address as your own when making purchases, which was as a work-related expense response. For fellow Canadian readers, the Canadian Radio Shack corporate address is:

279 Bayswater Drive
Brentwood, Ontario, Canada
L4S1A4W3
Tel: 705 728 6211

ProchBrou

Dear 2600:

In your previous response to a letter from "gr" in 18.2, I would like to pass this link on to the community: www.informationweek.com/contributors/1978/09/11/181800010. It details the entire DoCS519 file to the date of its publishing (July 16, 2001) in a very easy to read manner that is suitable for even the non-technical user who have no previous knowledge of the case. Although it may be a bit long (at today's attention depleted masses, it's not best), there come some errors.

Injustices

Dear 2600:

In regards to the letter by "SelfOut" in 18.1 about employers at Target not allowing you to use the Kodak Longs processor on a single-strike printout. I've worked at a copy center and while some printers are copyrighted (i.e., not reproduction without permission of the copyright holder), generally when a printer was involved, we worked the other way. I'm really not sure whether this is just a violation of working the law or how far reproduction under such circumstances, and it's probably not too far to get an agreement for your purposes, but my advice for those wanting this information is the former would be to take it to a suit where every court and opinion in the circumstances. You may want to inform the copy center of the financial harm doing the service as well, just as a good last-ditch. Most of the time, there will be someone who works there who will sympathize and allow you your last use.

Anonymous Partyseven

If a party gambler that never got arrested being subjected to this in the first place.

Dear 2600:

I was flipping through the channels on my TV and an unrec public access station there was a lecture on the Secret Service. They were going through all the algorithms and having people say what they do, as well as showing on some random issues. They got an interesting call phone 110's and state card fraud, and they showed a person at the other end of your website. It lasted after 30 seconds - they must have had a 24/7 or something. I just thought you'd find that interesting.

lurback

It's ironic that they sometimes announce an action such activities when we're always being quite vocal in our opposition to them. It's also ironic that a "public access" channel is being used for more government propaganda.

Dear 2600:

This is an excerpt from the July 2001 issue of *Sovereign America*:

"July 1981. The inexplicable counter-revolution and resurgence of the Tariffed customs authorities was accordingly shown by the prohibition of the importation of goods into the country. The reason advanced by the authorities was that in the event of sanctions violations, executed by the Department of Treasury, it would be impossible to obtain any other by which the operator of the machines could be traced. A large sum of \$200 million was being lying in the eastern house at the time the above law was passed and will have to be returned."

Who would have thought that after 180 years of alleged independence, we'd find ourselves in a similar situation as the country of Turkey was 150 years ago? "Digital Millennium Act" or "Ancient Turkish Revenge"?

MOXNETBOX

Compromising Internet Appliances

by Alex Hightmire

With today's exorbitant and today's concentration in the Internet has become larger than any other of mankind's creations. And everyone wants to be on it. People are rushing out to buy computers for the sole purpose of "getting on the net." With this bursting of wired technology and international networking, common everyday devices are now being made with interfaces to work through the Internet. With these new applications comes the inevitable security risks that come with every system on the net.

For example, there are several exercise devices that can be connected to the Internet, thus allowing the user to have a virtual trainer online guiding them and controlling their device. These are unannounced workouts that people can run through this company's website, www.fitness.com.

Web servers have been known to have exploits, allowing attackers to gain access to the system and permitting them to change any file on the server, including the graphics files that are used to control the exercise equipment during unattended workouts. If an attacker was to alter these workouts to force the runner to keep up a pace of 15 mph at a 20 percent incline, thousands of 50 year olds across the nation would either have a heart attack or fall off the spreading treadmill and hurt themselves.

Another fine example of a device that could be compromised is that of i-ready sexual devices. One such company, at www.safesites-plus.com/gasgas/SSE_Convertchart sells a device that attaches to your monitor. The box reads two parallel boxes that range from black to white. The intensity of redness controls the intensity of the vibration/rotation etc. All the attacker would have to do is replace the adjustable java applet with an saturated gif that alternates the extremes (black and white) which

would cause the devices to switch between off and high speed quickly, possibly burning out the device, but definitely annoying or harassing the user.

A final example is that of Internet appliances meant to reside in the kitchen of the house, allowing the user to listen to streaming music, browse sites (providing a recipe or two), watch DVDs, and monitor other appliances in the kitchen. The last option is the most vulnerable. At this time I believe it can only monitor the devices, but if an attacker broke into the appliances, they could possibly modify the software that monitors and calibrate it incorrectly, thus causing the turkey that is supposed to be finished cooking in one hour to remain in the oven for three hours before the user is alerted that it is done. Of course, freshness could ensue also.

These are just several of the existing devices that today you can expect to see more and more of these "Internet ready" appliances appearing in people's homes. Manufacturers of these appliances will face a whole new horror as consumers begin to lawsuit for loss of life, limb, or property due to a device being compromised. *Guest to Kismetville, Lord Manticore, Avatar, Error Blind, Hydrology, Kryptical, The Trance, Trancer and Mervin.*



An Introduction to ARP Spoofing

by Sean Whalen
arpspoof@gnx.net

This article deals with the subject of ARP spoofing. ARP spoofing is a method of exploiting the laxation of IP and Ethernet protocols. It is only applicable to Ethernet networks running IP.

Anyone with basic networking experience can understand key points of the subject. Knowledge of the TCP/IP reference model is vital to full understanding, as is a familiarity with the operation of switches and non-switched networks. Some background will be presented in the "Introduction" section, but experienced readers may wish to skip to "Operational".

Introduction

A computer connected to an IP/Ethernet LAN has two addresses. One is the address of the network card, called the MAC address. The MAC, in theory, is a globally unique and unchangeable address which is stored on the network card itself. MAC addresses are necessary so that the Ethernet protocol can send data back and forth, independent of whatever application protocols are used on top of it. Ethernet builds "frames" of data, consisting of 1500 byte (batches) each frame has an Ethernet header containing the MAC address of the source and the destination computer.

The second address is the IP address. IP is a protocol used by applications, independent of whatever network technology operates underneath it. Each computer on a network must have a unique IP address to communicate. IP addresses are virtual and are assigned via software.

IP and Ethernet must work together. IP communicates by constructing "packets" which are similar to frames, but have a different structure. These packets cannot be delivered without the data link layer. In our case they are delivered by Ethernet, which splits the packets into frames, adds an Ethernet header for delivery, and sends them down the cable to the switch. The switch then decides which port to send the frame to, by comparing the destination address of the frame to an internal table which maps port numbers to

MAC addresses.

When an Ethernet frame is constructed, it must be built from an IP packet. However, at the time of construction, Ethernet has no idea what the MAC address of the destination machine is, which it needs to create an Ethernet header. The only information it has available is the destination IP from the packet's header. There must be a way for the Ethernet protocol to find the MAC address of the destination machine, given a destination IP.

This is where ARP, the Address Resolution Protocol, comes in.

Operation

ARP operates by sending out "ARP request" packets. An ARP request asks the question, "Is your IP address x.x.x.x? If so, send your MAC back to me." These packets are broadcast to all computers on the LAN, even on a switched network. Each computer examines the ARP request, checks if it is currently assigned the specified IP, and sends an ARP reply containing its MAC address.

To minimize the number of ARP requests being broadcast, operating systems keep a cache of ARP replies. When a computer receives an ARP reply, it will update its ARP cache with the new IP/MAC association. An ARP is a stateless protocol, most operating systems will update their cache if a reply is received, regardless of whether they have sent out an actual request.

ARP spoofing involves constructing forged ARP replies. By sending forged ARP replies, a target computer could be convinced to send frames destined for computer A, to instead go to computer B. When done properly, computer A will have no idea that this redirection took place. The process of updating a target computer's ARP cache with a forged entry is referred to as "poisoning."

Attacks

Sniffing:

Switches determine which frames go to which ports by comparing the destination MAC on a frame against a table. This table contains a list of ports and the attached MAC address. The



table is built when the switch is powered on, by examining the source MAC from the first frame transmitted on each port.

Network cards can enter a state called "promiscuous mode" where they are allowed to examine frames that are destined for MAC addresses other than their own. On switched networks this is not a concern, because the switch routes frames based on the table described above. This prevents sniffing of other people's frames. However, using ARP spoofing, there are several ways that sniffing can be performed on a switched network.

A "man-in-the-middle" attack is one of these. What a MITM is performed, a malicious user inserts his computer between the communications path of two target computers. Sniffing can then be performed. The malicious computer will forward frames between the two target computers so communications are not interrupted. The attack is performed as follows (where X is the attacking computer, and T1 and T2 are targets):

X poisons the ARP cache of T1 and T2.

T1 associates T2's IP with X's MAC.

T2 associates T1's IP with X's MAC.

All of T1 and T2's IP traffic will then go to X first instead of directly to each other.

This is extremely potent when we consider that not only can computers be poisoned, but e-mail servers as well. All Internet traffic for a host could be intercepted with this method by poisoning a MITM on a target computer and the LAN's router.

Another method of sniffing on a switched network is MAC flooding. By sending spoofed ARP replies to a switch as an extremely rapid rate, the switch's port/MAC table will overflow. Besides early broadcast, but some switches will revert to broadcast mode at this point. Sniffing can then be performed.

Broadcasting:

Frames can be broadcast to the entire network by setting the destination address to FF:FF:FF:FF:FF:FF, also known as the broadcast MAC. By sweeping a network with spoofed ARP replies which set the MAC of the access gateway to the broadcast address, all external bound data will be broadcast, creating sniffing.

If a host were to listen for ARP requests and generate a reply containing the broadcast address, potentially capturing amounts of data could be broadcast on large networks.

DoS:

Updating ARP caches with non-existent MAC addresses will cause frames to be dropped. These could be sent out in a sweeping fashion to all clients on the network in order to cause a Denial of Service attack. This is also a side effect of port-MIM attacks, since targeted computers will continue to send frames to the attacker's MAC address even after they remove themselves from the communication path. To perform a clean MITM attack, the target computers would have to have the original ARP entries removed by the attacking computer.

Hijacking:

Connection hijacking allows an attacker to take control of a connection between two computers, using methods similar to the MITM attack. This transfer of control can result in any type of session being transferred. For example, an attacker could take control of a telnet session after a target computer has logged in to a remote computer as administrator.

Cloning:

MAC addresses were intended to be globally unique identifiers for each network interface produced. They were to be burned into the ROM of each interface, and not be changed. Today, however, MAC addresses are easily changed. Many users can even change their MAC without special software, using a simple parameter to "tkconfig", the interface configuration program for the OS.

An attacker could thus a target computer, then assign themselves the IP and MAC of the target computer, receiving all frames intended for the target.

Tools

ARPWatch implements packet-sniffing/monitoring. ARPWatch is a command line tool for UNIX which creates spoofed ARP replies. Users can specify the source and destination IP/MAC addresses.

Ettercap *http://ettercap.sourceforge.net*

Ettercap is a powerful UNIX program capable of a man-in-the-middle attack. It can be used by "sniffing traffic". All operations are automated, and the target computers are chosen from a sniffable list of hosts detected on the LAN.

Ettercap can perform four methods of sniffing: IP/MAC, ARP and Public ARP. It also automates the following procedures:
Injecting characters into connections.
Sniffing encrypted SSH sessions.
Password collision.

OS fingerprinting.
Connection hijacking.

Parasites:

Parasite is a daemon which watches a LAN for ARP requests and automatically sends spoofed ARP replies. This places the attacking computer as the MITM for any computer that broadcasts an ARP request. Essentially, this results in a LAN-wide MITM attack and all data on the switch can be sniffed.

Parasite does not do a proper cleanup when stopped. This results in a DOS of all poisoned computers because their ARP caches are pointing to a MAC address that is no longer forwarding their frames. Poisoned ARP entries must expire before normal operation can resume.

Defenses:

There is no universal defense against ARP spoofing. In fact, the only possible defense is the use of static (non-changing) ARP entries. Since static entries cannot be updated, spoofed ARP replies are ignored. To prevent spoofing, the ARP tables would have to have a static entry for each machine on the network. The overhead in deploying these tables, as well as keeping them up to date, is not practical for most LANs. Also of note is the behavior of static routes under Windows. Tests found that Windows will accept spoofed ARP replies and updates the static entry with the forged MAC, subverting the purpose of static routes.

MAC cloning can be prevented by a feature found on high end switches called Port Security (also known as Port Binding or MAC Binding). Port Security prevents changes to the MAC entries of a switch, unless manually performed by a network admin. It is not suitable for large networks, or networks using DHCP. Port Security does not prevent ARP spoofing.

Aside from these two methods, the only remaining defense is detection. Arpwatch is a free UNIX program which listens for ARP replies on a network. It will build a table of IP/MAC associations and store them in a file. When the MAC address associated with an IP changes (formed to as a flip-flop), an email is sent to an administrator.

Tests showed that running Parasite on a network caused a flood of flip flops. Leaving the MAC of the attacker present in Arpwatch's email. Ettercap caused several flip flops, but would be difficult to detect on a DHCP-enabled network where flip flops occur as regular intervals.

MAC cloning can be detected by using RARP (Reverse ARP). RARP requests the IP address of a known MAC address. Sending a RARP request for all MAC addresses on a network could determine if any computer is performing cloning, if multiple replies are received for a single MAC address.

If a MAC level is performed and the switch reverts to broadcast mode, a computer will have to enter promiscuous mode to examine the broadcast frames. Many methods exist for detecting machines in promiscuous mode. These can be found in the sniffing FAQ, at <http://www.robeygram.com/pktsniffing/>, without being in promiscuous mode since redirected frames will be routed to your MAC.

It is important to remember that operating systems have their own TCP/IP stacks and Ethernet cards have their own drivers, each with their own quirks. Even different versions of the same operating system have variations in behavior. Schirer is unique in its treatment of ARP replies. Not only does it accept ARP updates after a timeout period. To poison the cache of a Solaris box, an attacker would have to DOS the second target machine in order to avoid a race condition when the timeout period. This DOS may be detected if the network has an Intrusion Detection System in place.

Cloning:

ARP spoofing is one of several vulnerabilities which exist in modern networking protocols, which allow a knowledgeable individual free reign over a network. IP spoofing, TCP sequence prediction, and ICMP offsets are just a few examples of other current weaknesses in these protocols. It is unlikely that these problems will be addressed until they are solved on a wide enough scale to force a change in the standards quo. The problem is poised to grow as broadband Metropolitan Area Networks are implemented using Ethernet as the protocol of choice.

Information in this article was heavily referenced by the *Metasploit* and *Parasite* programs. Proof of concept tests were performed with the tools mentioned here, against Linux, Windows NT, and Windows 2000 machines.

OFFSET HACKING or How I got banned from EverQuest

by Darbyesh

Client hacking, the process of finding out which offsets affect what using a hex editing utility like WinPacker or Soft Ice, has been around for a long time. Games like Diablo have been quite well known for it. However, they never seemed to ban players for it. As a matter of fact with Diablo 2, they made everything server side to stop that process. But in the spirit of good fun they still have their open bootloader which still can be modified and offers the option of play in areas who want to play that way.

In all honesty probably every online game that's out there gets hacked. Why can't it be fun. Typically it starts off with just people the same way. They play the game a long time, then get bored. Then they start to try and figure out ways to hack the game. Think of it - you don't have to worry about going to jail and you still have the pleasure of having the systems. Considering most of the people who do it are power gamers and put enough money into their game computers, pockets, the companies usually don't care that much about it.

Then along comes Verant Inc. with their all out client Everquest game or Everquest as most like to call it. First off let me state this company has probably had the worse sort of customer support right from the start. Their so called guides are nearly useless and usually will say they need to refer you to a nice guy who is never met. Not to mention all the bugs they still have. I fixed 4 bugs on after two years you would think they'd get it right.

Obviously you being all these elements together and of course now after hacking Everquest seems like a mighty invasive proposition. And you may even think they don't care since they obviously don't care enough about good customer service. Wrecking sites of their authentic about any applications came when Dan Ziegler came up with a macro utility to make game play better. It was called EQ Macros. He was stopped in his tracks as you can see from this quote from his website: "EQ Macros is temporarily banned - it is not being sold or distributed. Also, send the CEO of Verant, send me an email and request that I stop work on EQ Macros. I requested asking him to consider developing a 3rd party developer support program, that Original EQ Mac program, so that we could work together on improving the EQ gaming experience. When received communications from Ver-

ant I haven't asked for the name of the guy who I wrote this.

What's also interesting is that Verant wanted to be able to scan your PC for third party apps. They changed their mind after users protested. It would have explained why to allow Verant to upload any data that could "interfere" with the proper operation of EverQuest."

Now the thing to keep in mind is that the offset hacking is going on at the other end of the game which sits on your hard drive. So what Verant is trying to do is that you can't sniff it out at something that is on the PC that you bought. Also consider the fact that you bought the software as well as paying a \$10 a month subscription fee.

Recently Verant and their Xmas signed banned over 500 accounts for hacking. Now when people asked for proof they received some nice Latin letters. The human contact was not offered. Many users were first greeted with this email:

"It is my regretful duty to inform you that your Everquest account... has been banned for violating our Agreement. Bank of Canada and our OpenQuest User Agreement and Support License to which you have affirmatively agreed to abide by each and every time you play EverQuest. The use of a third party program to alter your gameplay is not permitted and as such has warranted the removal of your access to the game.

"If you have any questions or concerns regarding this action please feel free to contact our customer support team at support@verant.com. As a result of this action the registered credit card will no longer be valid for the EverQuest subscription fee.

"We thank you for your past patronage."

Now granted we all know the risks and we paid the price. But what I find interesting is that they will not offer proof to back up their claim that you were hacking. Which leads me to believe that they are doing some kind of client side scanning. I do know some innocent bystanders did get banned. If Verant knows this is not going to offer up proof to have they know your hacking. I think they should remove your account. Or is the real truth that they are scanning your PC which is an invasion of your privacy?

I think the main bannings with no offering of evidence is at most the same as Kevin Winfield's case. What's the secret coming to when some gaming companies can get away with this stuff? Seeing the customer it always right obviously does not apply to Verant.

One thing the company doesn't realize is that while it's only \$10 a month, the time put into building up those characters was a lot more than that. The fact that Verant is not showing how they caught you and just answering with a form letter is BS.

I will end this with a post that John Smedley himself put on the Hackquest board

This message is addressed to those of you that are examples of no hack EverQuest.

Read what messages on this board. YOU' can judge. How will you feel that a large number of people are being banned today.



by Lucky225

Lucky225@verantforums.com

The invisible box will make it so that when you pick up a phone on your phone line any of these in-cue lights that tell if an extension phone is picked up won't light.

Theory

The theory is based off the same principles as the infamous black box that used a 1.8k resistor to keep the phone line at 50v when you pick up. If successfully still works, but because of modern switching the voltage is cut off from the party calling you, and the phone company doesn't allow a voice connection anymore until your phone goes off hook and there's super-tension. The invisible box works by using high resistance to keep the voltage at about 30 volts. This is accomplished by placing a resistor of about 470ohms in series with your phone. The phone is approximately 215ohms and draws 28ma of current, which means when your phone is off hook there are approximately 6 volts on the phone line. When you place the resistor in series with the phone line, there is a total resistance of 685ohms. Using ohm's law, 685 ohms times 28ma gets you 19.2 volts. So the resistor keeps the phone line at about 20 volts, and most to use lights only go off when there are about 15 volts or less on the phone line.

Construction

You will need a phone cord and a 470ohm resistor (yellow, purple, brown). You can get the resistor in a five pack at Radio Shack for \$0.49. It wouldn't hurt to have some wire strippers and possibly electrical tape or solder. Strip the phone cord in the middle. Don't cut the middle jacks off.

You have been logging things on the server. You some time and will continue to do so in the future. If you back, you will be caught, and you will be banned as well as any.

Regards,

John Smedley,

Chief Operating Officer

South Pacific Entertainment

By the way, the thing he claims about logging is BS. If they were logging as they claimed, my friends and I have been banned as well. But they were forbidden enough to be out of town during that week.

You'll see four or two wires, usually black, red, green, and yellow. Don't worry about the black and yellow wires. In fact, cut them off as they'll get really weird. Leave the green wire alone. That's the positive wire, and since extension phones from negative to positive and we're trying to oppose current so the voltage won't drop, we leave it alone. Finally, cut the red wire (that's the negative) in half and strip both ends. You're going to insert the resistor here.

Conclusion

That's it. Pretty simple huh? You might be thinking that maybe there is no real use for this because all it does is make it so that in-use light doesn't light when you pick up the phone. But think of the possibilities. You could go badge boxing with this box and it might save you if the person you're talking to of this so in-use light and they always look at it to see if their kid is on the phone. Because of your trusty invisible box hooked up to the phone line, that light never comes on and they never pick up to yell at who they think is their kid. Personally, I use it when I'm talking on my phone line but want to use my main line to go on the internet. My room is always cluttered that damn in-use light and yelling at me. "You're on the internet with my phone line! Get off now!" But that's how you'll never know! The sad thing is that this even bypasses those lame \$200 phone tap devices you always see on TV.

Cover: *Ray*, my beloved! *Spencer*, *Kevin*, *BigB0000*, *Wagon*, *Clancy*, *Montal Angel*, *Kennedy*, *gavin*, *cm*, *Quynha*, *Emmett*, and anyone else I left out!

Deal with this responsibly.

Avatar

Dear 2600:

First of all, thanks for helping us fill in gaps for us on the subject of freedom. I've a couple hundred and don't pretend to understand all that bawling can talk, but the recent reversal of the Microsoft deal, order just begs the question: since, and now, Xerox Mexico bases out its exchange for their legal fees being paid. Having me, our government is sending out of court just to recon legal expenses? What's driving this fucking thing anyway?

(epdsh999)

Dear 2600:

I hear the final build of Windows XP is going to be 2850. Congratulations! I think not.

Dawson

It's not the renaming the price.

Just Plain Evil

Dear 2600:

If you think having your face scanned in public using cameras and computers to see if you are a "criminal" is bad, hold on. It's getting worse.

A public face-scanning system is already up and running in Tampa, Florida. Cameras are mounted in high crime neighborhoods, monitoring pedestrians in the streets. Using software called "FaceIt" by Vericon, US, snapshots are stamped against a database of 300,000 people that includes mugshots, fingerprints and reds wanted on any criminal charge. The police are dispatched when the software makes a match.

The Colorado DMV is installing new software to make it easier for the government to find you. When you have your picture taken for a driver's license, special 3D mapping software will be used to create a "facial" file, not unlike a fingerprint. The file contains information which identifies facial characteristics unique to that individual. "The file then becomes part of a database shared by government agencies. This way, if there's a camera the government can monitor (airways, post office, street corner, etc.) they can get you."

I am told by a friend in law enforcement that they are testing a mobile version of the Vericon's "FaceIt" system. The idea is that if you're riding in a traffic light and a cop pulls up next to you, sensors in the police car will automatically scan your face to see if you've been wanted. A laptop in the police car will show the officer if the software comes up with a match. It won't matter when faced of crime you committed. The system won't discriminate between burglar, robbers and parking ticket violators.

In Ontario, California police are testing a portable, wireless fingerprinting device by Motorola called FINGER. Upon demand, an officer asks you to place your finger on the device which then searches a database for an identity match. The device has a small video screen which also displays information about your identity and any outstanding warrants. If no matches occur, a built-in camera takes your picture and records

your fingerprint, picture, and personal information into the database.

Police in Colorado are testing a handheld reader device that can see through your clothes. The device allows officers to scan you for any concealed items at a distance. When used to scan a crowd it displays pictures on a built-in video screen.

All of this sounds like serious privacy issues. For more, see: <http://itj's.org>.

The ACLU is protesting all of the above with little success. Face scanning is "a virtual liturgy that struts of Big Brother by randomly monitoring people without their consent. And that technology does give law enforcement specialized powers - powers that go well beyond what would be provided by human senses," says Diane Steinhardt, associate director of the ACLU. "It allows police officers to engage in intrusive searches while you sleep."

Speed Bearer

Let's see, Government? Government? Member? Or individual people? How do we know the decision you just want a change.

Just Plain Stupid

Dear 2600:

This Coder@C worm program is quite interesting. I am watching it work on my personal computer, simply to see only two files for friends of mine. On August 16, my logs showed I had queries containing the script file used for the MS buffer overflow bug that Coder@C threw on. I took the time to look up all those IP addresses, then spotted MS servers, and of all the IP's, single host, egypt.400000000, many of which no longer work, were being a Microsoft cable connection. Most are corporate servers, obviously, as they are the only people who would pay for such a shoddy product by Microsoft. It is amazing how a 16-year old maintenance person works with better records than such giant mega corporations. Microsoft: Beware! Beware!

Lucius

AM Idea

Dear 2600:

I think it is time to ban the likes on the new copyright laws and use them to our own advantage and show the world how stupid they are. How about someone our there, even me, create a virus, get a copyright on it, and "accidentally" release it? Then when all the anti-virus software companies come out with an exception for the virus and having reverse engineered it, they get sued for violating the digital copyright law. And we all just for massive arrests. Maybe then could be the starting point to show how dumb the law really is.

David

With the general public, it's far less than without a second thought.

Net Jacking for Complete Idiots

by David Overton of the Doc

The best big thing in hacking these days is wireless 802.11 networking. The reason for that is that the hardware is cheap and open networks are abundant.

Wireless networks are popping up all over the place. From corporate offices to trade shows, conference halls, libraries, coffee shops, parks and personal residences.

In the corporate environment the majority of wireless LANs (WLANs) are connected to the internal backbone of the company. In other words, their corporate firewalls, thus unbeknownst, giving everyone within a two-block radius full unrestricted access to the internal network and attached outgoing resources.

Not all networks are private networks - there are many that are intended to be accessible to the public to protect business. For example, there are many coffee shops that are offering free Internet access in hopes that people will spend more time drinking coffee at their establishment. Hotels are providing wireless access as a perk to attract guests, which is often cheaper than wiring all the rooms with cable for Internet access.

In Seattle, San Francisco, and other cities, there are groups and organizations that are setting up WLANs for free use in their neighborhood in a philanthropic manner.

What is WLAN802.11?

802.11 is a standard for WLANs developed by the Institute of Electrical and Electronics Engineers (IEEE). The standard deals with network association, data transfer, authentication, and priority.

802.11 is the first draft of the protocol specifying transmission speeds of one and two megabits a second. The 802.11b specification describes a later update to the protocol for eleven megabit rates. (802.11a is a specification for 54 megabit rates but is not ready for prime time.)

The 802.11 WLAN protocol specifies the internal layer of the OSI network model (physical) from which other protocols such as TCP/IP, FTP, SSH, etc. build on.

On a traditional wired network, physical connectivity defines the physical description for use of layer two switching and VLANs. This assembly of nodes and words is physical. Physical networks, however, such as WLANs, are not physical connections to something.

Instead, the way networks are implemented is through means called SSIDs. To connect to a wireless network, all you need is the network name and to be within radio range of the wireless bridge. The SSID was never meant to provide that security but sadly that is how it is currently being used in current deployments.

The 802.11 protocol supports a layer three encryption method called WEP (wired equivalent privacy). WEP is a simple algorithm based on RC4's block hashing scheme. Recent analysis has shown WEP encryption

to be inadequate. The details of WEP and its weaknesses are beyond the scope of this article and will be the topic of a future article.

Hardware

For around \$50 to \$100 you can get a good PCMCIA card. I have seen older cards on ebay.com and on sale over the \$100-40 \$15. With this card and a standard laptop you can be working down the street or sitting in the park with free Internet access. I recommend the Lincant cards for their features: external settings options and cross platform support.

If you plan on setting 802.11 (example, I suggest a card based on the Prism chip set).

Software

All popular operating systems have the driver support for this cool popular card. (Linux/OS/UNIX, OS/2, Windows, Mac OS, etc.) For the examples used in this document we will use Microsoft Windows since it is the easiest to set up and install.

Locating a WLAN

Locating a network is easy once you get the hang of it. For simplicity, I will explain how to do this under Windows with a Lincant/OS/UNIX card and software. Once you have your card installed run the Client Manager software and set your SSID to "ANY". From the "Advanced" menu of Client Manager select "Site Machine". A window should open listing all the local WLANs that are available from where you are standing. I recommend doing this first near a known network, then as you move around click on the "Scan Now" button to refresh the view (see Figure 1 for an example output from the Client Manager in San Francisco).

If you do not have a Lincant/OS/UNIX card you can manually search for common SSIDs. All cards come with no configuration to quickly change the SSID you are using. Simply prepare in the file or see more information SSIDs (see Figure 2 for a list of common SSIDs) and cycle through them, or just walk around until you get a hit.

A more effective method is to sniff the 802.11 frames, and look for known SSIDs. This will require special software, often that is part of the scope of this introductory tutorial will be the topic of a future article.

What's Next?

At this point you can put "enough" to get a DHCP lease on an IP. After you get an IP address, you are on their network. You will be able to access the Internet (relatively anonymously). When you get on, click "Network Neighborhood". You should be able to see the other hosts and shared resources available on the network (remember - if you can see them, they can see you).

If you have a wireless path to NetXray or direct in scaled on your system, you will be able to see packets sent to other wireless hosts and broadcast packets.

